

# 061 - Final - Practice Problems

1 June 2011

1. Prove that  $n! > n^2$  for all integers  $n \geq 4$ .

**Solution** Proof by induction. Base case: when  $n = 4$ ,  $n! = 24$ , while  $n^2 = 16$ , so this checks out. Now, suppose that  $n > 4$  and that the statement is true for all  $k$  where  $4 \leq k < n$ . Then,

$$n! = n((n-1)!) > n((n-1)^2) = n^3 - 2n^2 + n = n^2(n-2) + n > n^2,$$

as desired.

2. Let  $X$  be a finite set with  $n$  elements. Determine, with proof, how many binary equivalence relations there are on  $X$ .

**Solution** A binary relation on  $X$  is just a subset of  $X \times X$ . The subsets of  $X \times X$  are the elements of the power set  $P(X \times X)$ . The set  $X \times X$  has  $n^2$  elements, so the set  $P(X \times X)$  has  $2^{(n^2)}$  elements. Therefore, there are  $2^{(n^2)}$  binary relations on  $X$ .

3. How many rearrangements of MATHEMATICS are there where the Ms are not next to each other?

**Solution** In general, there are a total of

$$\frac{11!}{2!2!2!}$$

rearrangements of MATHEMATICS. Let  $\Phi = \text{MM}$ . Then, there are

$$\frac{10!}{2!2!}$$

rearrangements of  $\Phi\text{ATHEATICS}$ . These correspond to the rearrangements of MATHEMATICS in which the Ms *are* next to each other. So, there are

$$\frac{11!}{2!2!2!} - \frac{10!}{2!2!}$$

rearrangements of MATHEMATICS where the Ms are not next to each other.

- 4.** Let's play Canasta! The deck consists of 2 standard packs of 52 cards, 13 in each of 4 suits. So, there are 2 of every card, but we can't tell the two copies apart. For example, there are 2 Aces of Hearts. How many different 5-card hands are there that contain only Hearts?

**Solution** First, suppose that the hand contains no duplicates; e.g., there are not 2 Aces of Hearts in the hand. Then, there are  $\binom{13}{5}$  such hands. Now, suppose that a single card is duplicated. There are 13 choices for the duplicated card, and  $\binom{12}{4}$  choices for the other cards. If 2 cards are duplicated, there are  $\binom{13}{2}$  choices for those cards and  $\binom{11}{1}$  choices for the other card. Therefore, there are

$$\binom{13}{5} + \binom{13}{1} \binom{12}{4} + \binom{13}{2} \binom{11}{1}$$

different flushes of Hearts.

- 5.** Let  $X = \{1, 2, 3, 4, 5\}$ . How many strings of length 1000 on  $X$  are there such that there are no substrings from  $\{1, 2\}$  of length more than 1.

**Solution** Let  $a_n$  be the number of string of length  $n$  on  $X$  such that there are no substring from  $\{1, 2\}$  of length more than 1. Then,  $a_0 = 1$  and  $a_1 = 5$ . We find a recursive formula for the  $a_n$ . Given any string  $t$  of length  $n - 1$  on  $X$  of the same type, the strings  $3t$ ,  $4t$ , and  $5t$  are all of the appropriate type. Similarly, given any string  $t$  of length  $n - 1$  on  $X$  of this type, the strings  $13t$ ,  $14t$ ,  $15t$ ,  $23t$ ,  $24t$ , and  $25t$  are of the correct type. Thus, we see that

$$a_n = 3a_{n-1} + 6a_{n-2}.$$

To solve this, we consider the equation  $t^2 - 3t - 6$ . Using the quadratic formula, this has solutions  $r_1 = \frac{3+\sqrt{33}}{2}$  and  $r_2 = \frac{3-\sqrt{33}}{2}$ . Solving the system of equations

$$\begin{aligned} a + b &= 1 \\ ar_1 + br_2 &= 5, \end{aligned}$$

we find that  $a = \frac{7}{\sqrt{33}}$  and  $b = 1 - \frac{7}{\sqrt{33}}$ . Therefore, there are

$$\frac{7}{\sqrt{33}} \left( \frac{3 + \sqrt{33}}{2} \right)^{1000} + \left( 1 - \frac{7}{\sqrt{33}} \right) \left( \frac{3 - \sqrt{33}}{2} \right)^{1000}$$

such strings.

- 6.** Prove that in any set of 51 positive integers less than 100, there are two whose sum is 100.

**Solution** Let  $a_1, \dots, a_{51}$  be 51 positive integers less than 100. Let  $b_n = 100 - a_n$ , for  $1 \leq n \leq 51$ . First, note that  $b_n = a_n$  if and only if  $a_n = 50$ . If some  $a_n$  is equal to 50, then discarding  $a_n$  and  $b_n$ , the rest of the numbers form 100 integers between 1 and 99. Thus, two of them are equal by the pigeonhole principle. So,  $a_k = b_j = 100 - a_j$  for some  $k \neq j$ . So, we're done. If no  $a_n$  is equal to 50 then the same argument works.

7. Show that if  $G$  is a simple graph, then either  $G$  or  $\overline{G}$  is connected.

**Solution** Assume that  $G$  is a simple disconnected graph. Let  $v_1, \dots, v_k$ ,  $k \geq 2$ , be a vertex from each connected component of  $G$ . This means that every vertex of  $G$  can be connected to exactly one of the  $v_i$ , and no  $v_i$  can be connected to any other. Let  $x$  and  $y$  be two vertices in the vertex set of  $G$ . We show that they are connected by a path in  $\overline{G}$ . First, if  $x$  and  $y$  are in different components in  $G$ , then there is actually an edge between them in  $\overline{G}$ , so they are certainly connected by a path in this case. Now, assume that  $x$  and  $y$  are in the same component of  $G$ , say the  $v_1$  component. Then, there is an edge  $e_1$  from  $x$  to  $v_2$  in  $\overline{G}$  and an edge  $e_2$  from  $y$  to  $v_2$  in  $\overline{G}$ . Thus, the path  $(x, e_1, v_2, e_2, y)$  in  $\overline{G}$ . Therefore, in this case too  $x$  and  $y$  are connected. Therefore,  $\overline{G}$  is connected.

8. Show that if  $G$  is a simple graph with at least two vertices, then there are two vertices in  $G$  with the same degree.

**Solution** Suppose that  $G$  has  $n$  vertices. Since  $G$  is simple, the degree of each vertex is between 0 and  $n - 1$ . If the graph is connected, then the degree of each vertex is between 1 and  $n - 1$ . By the pigeonhole principle, two vertices have the same degree. If the graph is not connected, there is no vertex of degree  $n - 1$ . Thus, the degree of each vertex is between 0 and  $n - 2$ . Again, by the pigeonhole principle, two vertices have the same degree.

9. Prove that every tree with at least two vertices is a bipartite graph.

**Solution** Choose a root for the tree  $T$ . Then, let  $X$  consist of the vertices of even level, and let  $Y$  be the vertices of odd level. Then,  $T$  is bipartite on  $X$  and  $Y$ .

10. Prove that the number of nonisomorphic binary trees with  $n$  vertices is the  $n$ th Catalan number.

**Solution** Denote by  $C_n$  this number. Then,  $C_0$  is 1. We can construct all isomorphism classes of binary trees with  $n$  vertices by choosing the number of vertices  $k$  of the left branch of the root together with a binary tree on  $k$  vertices together with a binary tree on  $n - k - 1$  vertices. Therefore,

$$C_n = \sum_{k=0}^{n-1} C_k C_{n-k-1}.$$

But, this is the same recurrence relation satisfied by the Catalan numbers with the same initial condition.